



MIDAS: Middlebox Discovery and Selection for On-Path Flow Processing

NFV workshop 2015

Published on Comsnets 2015

TNOVA

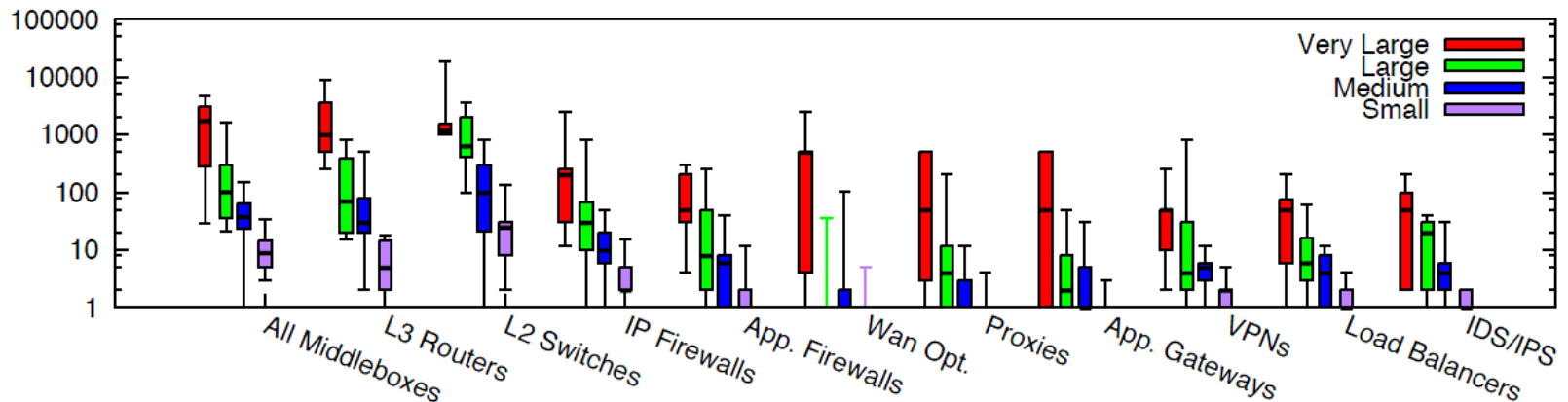


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- Today's enterprise relies on wide range of middleboxes :
 - packet filtering
 - proxies
 - load balancing
 - redundancy elimination
 - encryption
 -



J. Sherry et al., **Making Middleboxes Someone Else's Problem: Network Processing as a Cloud Service**, SIGCOMM 2012



- Today's Middleboxes limitations:
 - Specialized Hardware and functionality
 - High investment cost
 - Standalone device provisioned for peak loads
 - Inefficient resource utilization
 - Diverse management and configuration interfaces
 - High operation cost
 - Deployed closed to the edge
 - Single point of failure
 - Concentrate traffic at the edge
- NFV, replaces middleboxes with software NFs on:
 - To the cloud [APLOMB] (off-path)
 - To the network (on-path)

J. Sherry et al., **Making Middleboxes Someone Else's Problem: Network Processing as a Cloud Service**, SIGCOMM 2012



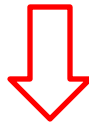
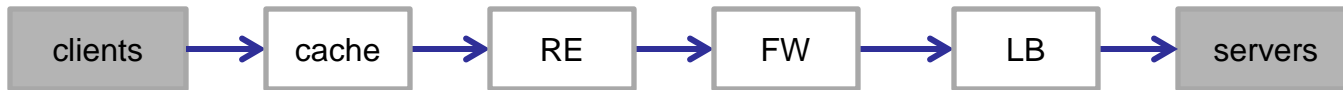
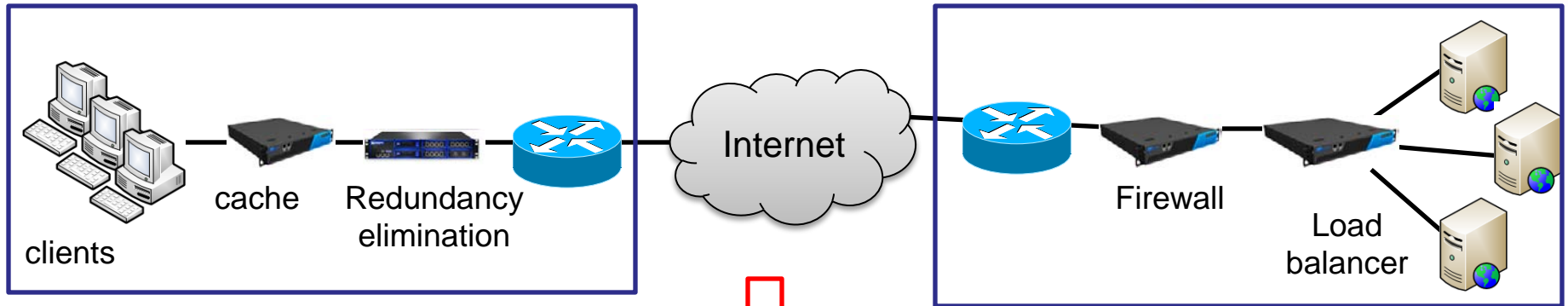
- Migration of middleboxes to the network:
 - Enable NFaaS
 - Pay-per-use model
 - Reduced CAPEX/OPEX and high flexibility and scalability
 - Empowering the “middle”
 - Bandwidth conservation
 - Redundancy elimination
 - Packet filtering for DoS mitigation

- Recent trends
 - Routers with programmable processors
 - Packet processing on commodity servers
 - Consolidated SW middleboxes [Flowstream, CoMB]
 - Micro-datacenter deployment by ISPs

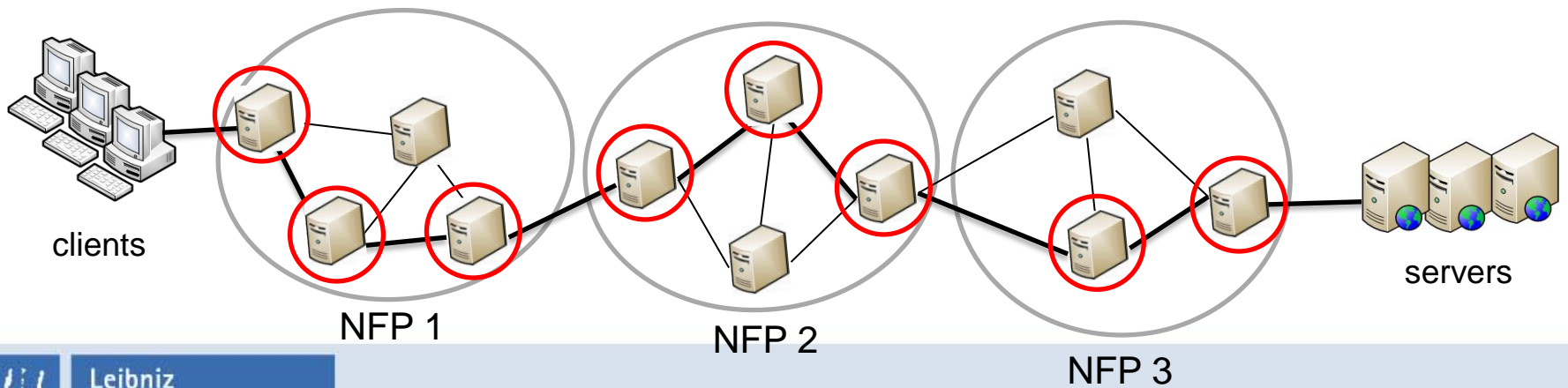
A. Greenhalgh et al, **Flow Processing and the Rise of Commodity Network Hardware**, CCR 2009

V. Sekar et al., **The Design and Implementation of a Consolidated Middlebox Architecture**, NSDI 2012

On-Path Flow Processing

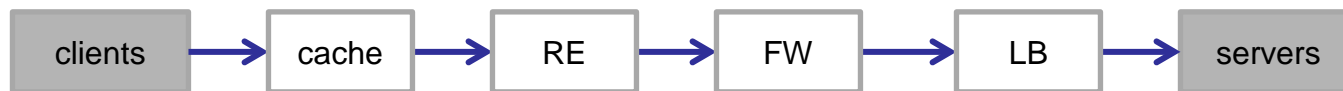


Middlebox ~~ab~~o





- Performance
 - High packet forwarding rates [RouteBricks, ClickOS]
 - Low processing setup delay
- Resource utilization efficiency
 - Load balancing
- Correctness
 - Network functions (NFs) should be embedded in the correct order
- Proximity
 - Some NFs should be closed to the source/destination

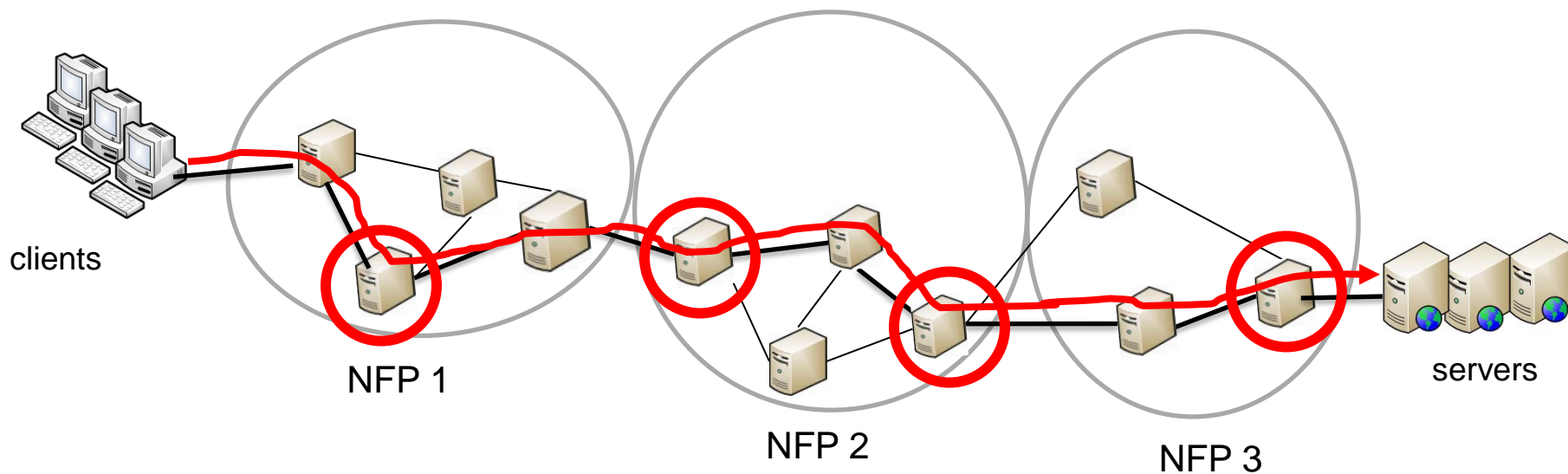
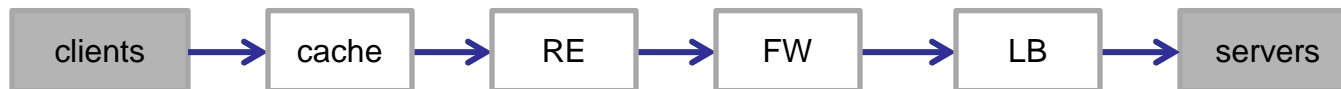


M. Dobrescu et al., **RouteBricks: Exploiting Parallelism to Scale Software Routers**, SOSP 2009

J. Martins et al., **ClickOS and the Art of Network Function Virtualization**, NSDI 2014



- Middlebox discovery
 - Path discovery and middlebox detection techniques (e.g., traceroute, tracebox) incur high delays
 - Signaling protocols (e.g., SIMCO) are designed for middlebox configuration
- Middlebox selection
 - NF location dependencies require large provider footprint (i.e., multiple NFPs)
 - NFP resource information disclosure policies



- Middleboxes pick up flows as they arrive
 - ✓ Performance
 - ✗ Trade-off between correctness and efficiency
- Need for processing setup coordination within and across NFPs



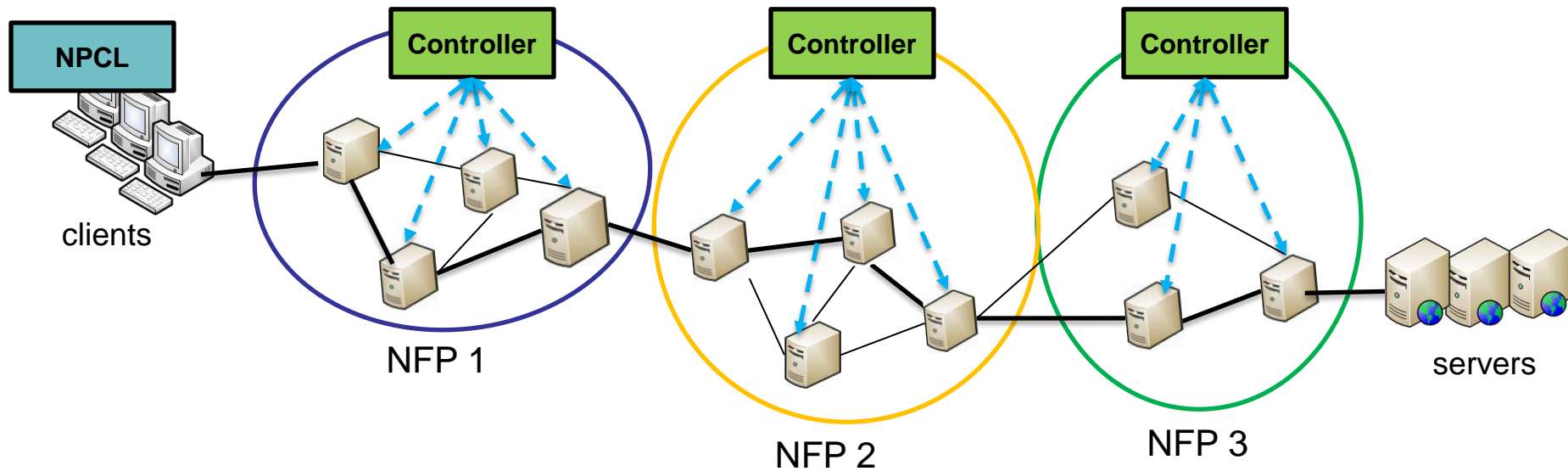
- MIDAS overview
- Middlebox discovery
 - Middlebox signaling
 - Middlebox controller chaining
- Middlebox selection
 - Privacy-preserving NFP assignment
 - Intra-provider middlebox selection
- Implementation
- Evaluation
- Conclusions

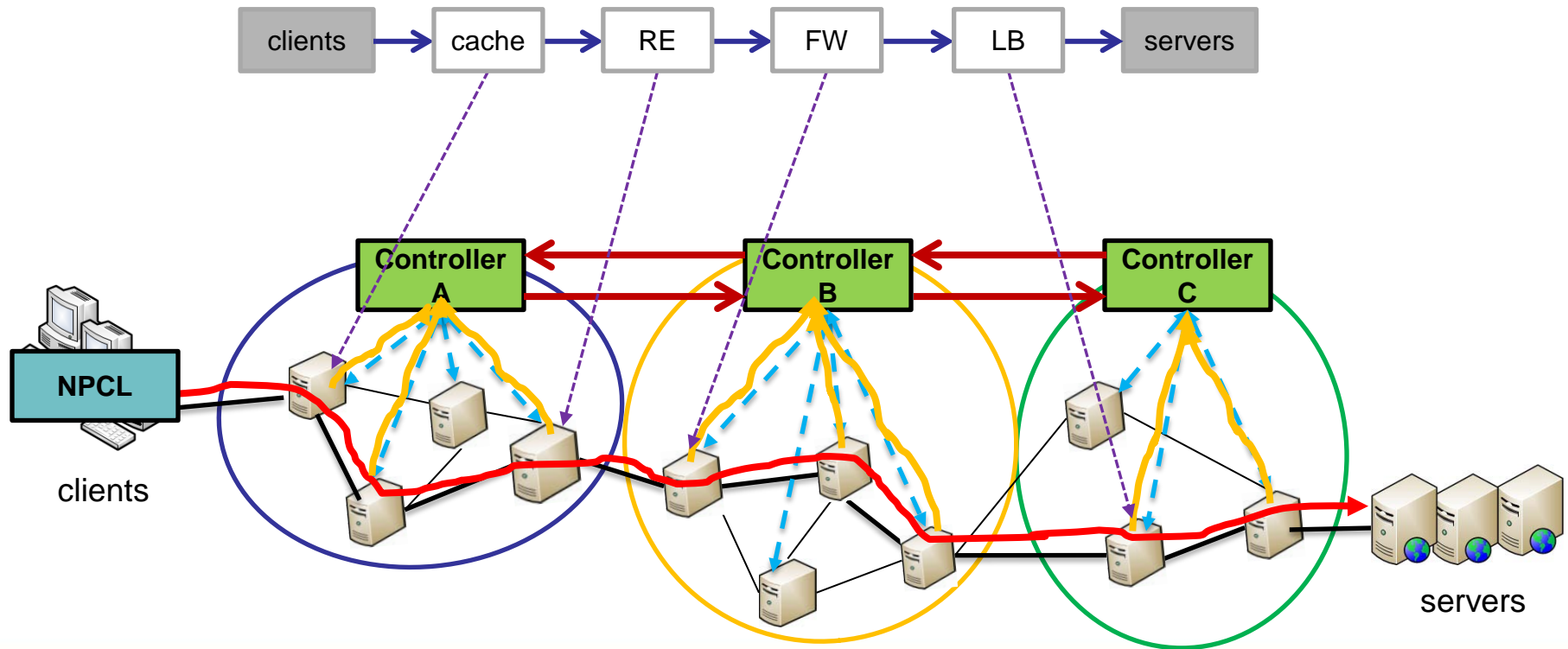


MIDAS Overview



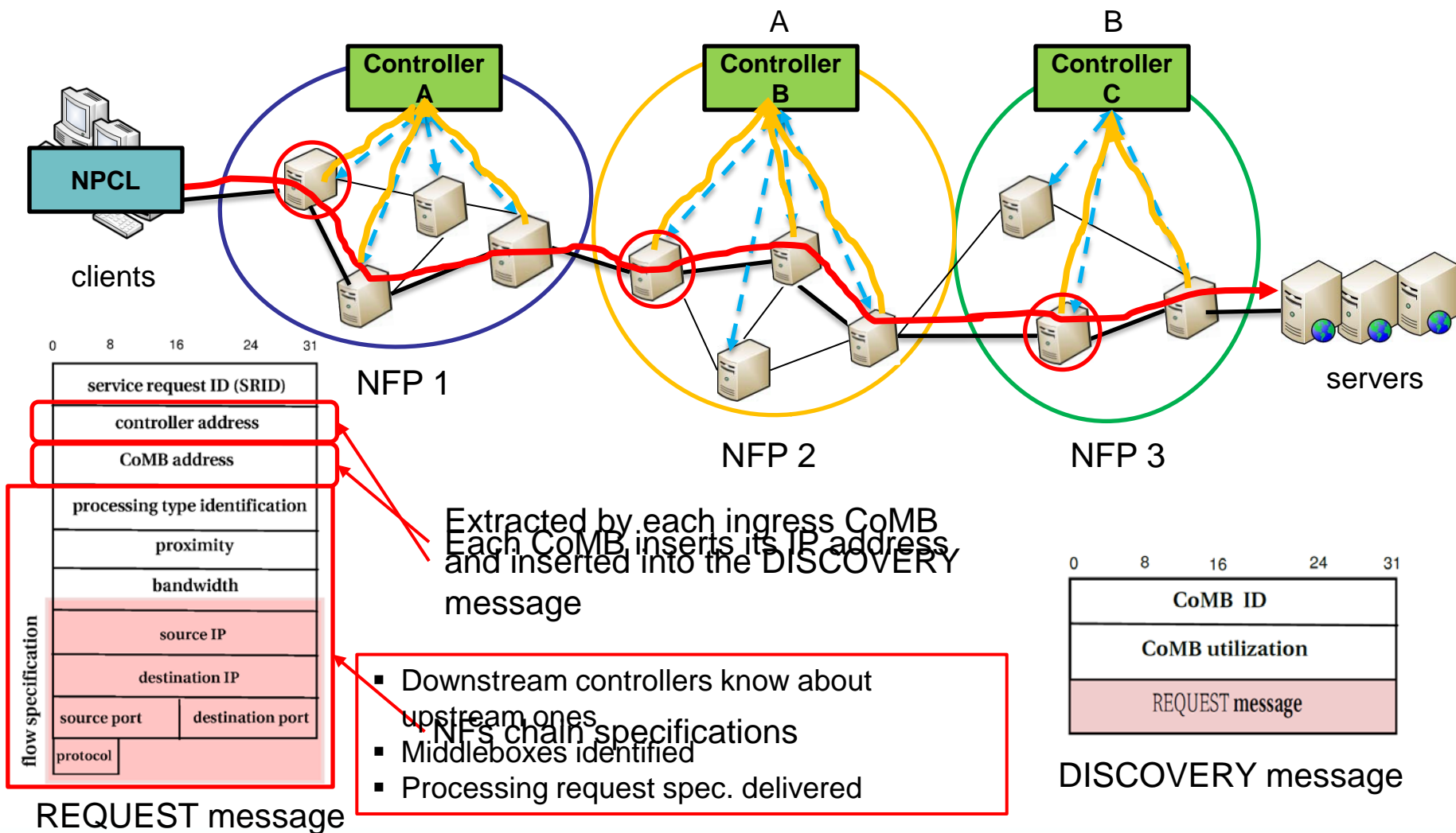
- Main components:
 - Consolidated middlebox (CoMB)
 - Centralized CoMB controller in each NFP
 - Network processing client (NPCL)

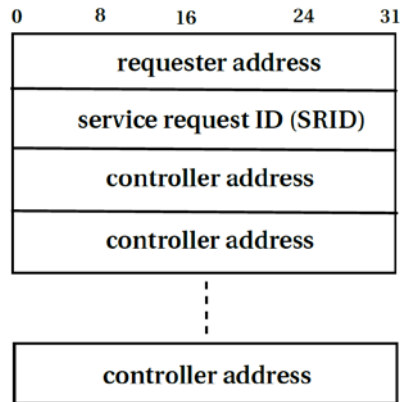
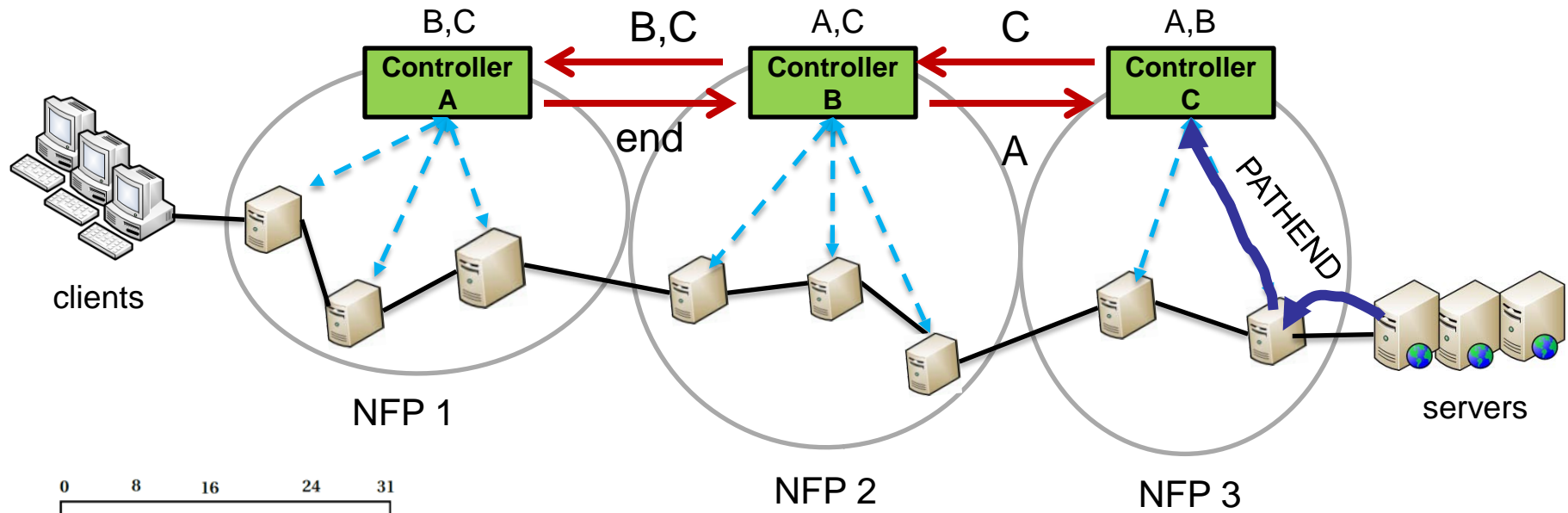






Middlebox Discovery





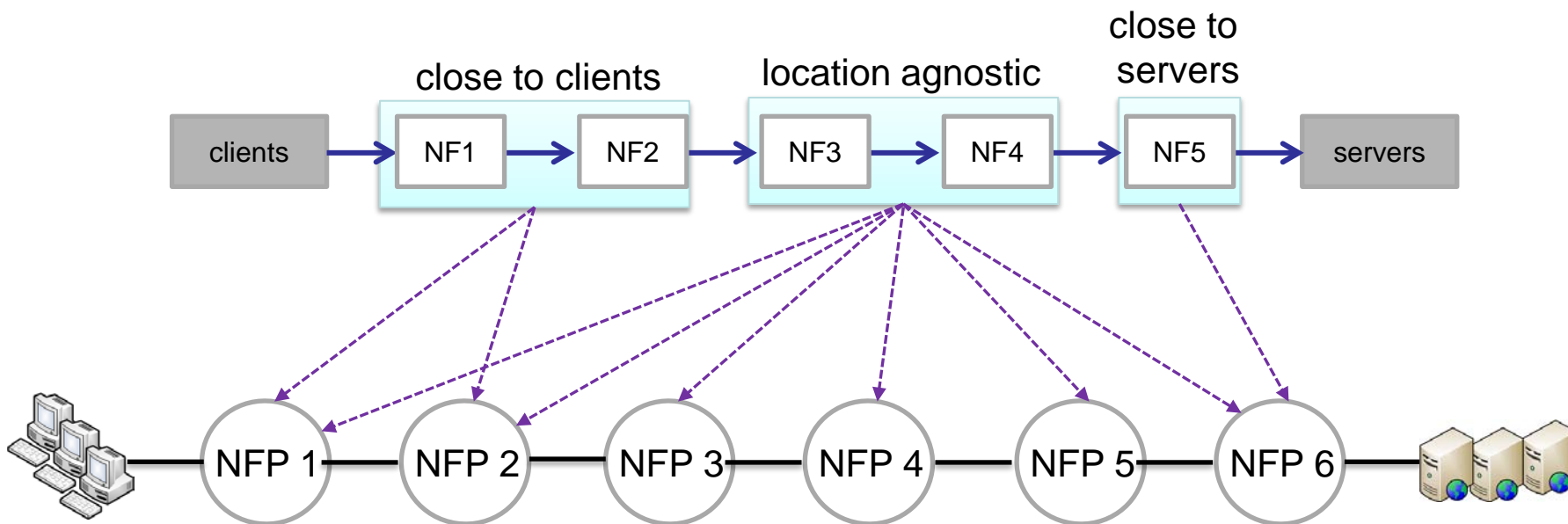
CONTROLLER message



Middlebox Selection



- Objective:
 - Minimize number of assigned NFPs
 - Maintain providers privacy
- Distributed approach:
 - Each NFP partitions the service chain based on NF location dependencies
 - Each NFP announces which segment it is willing to host
 - Privacy-preserving protocol (using MPC) to assign chain segments to the NFP with the lowest utilization





- Cryptographic protocol :
 - Different parties with private inputs to compute a function on their inputs:
 1. Input values stays private
 - Utilization of each NFP
 2. Result of the computation is correct
 - Compute the NFP with the lowest utilization
 3. Cheating parties won't learn information about the honest parties inputs

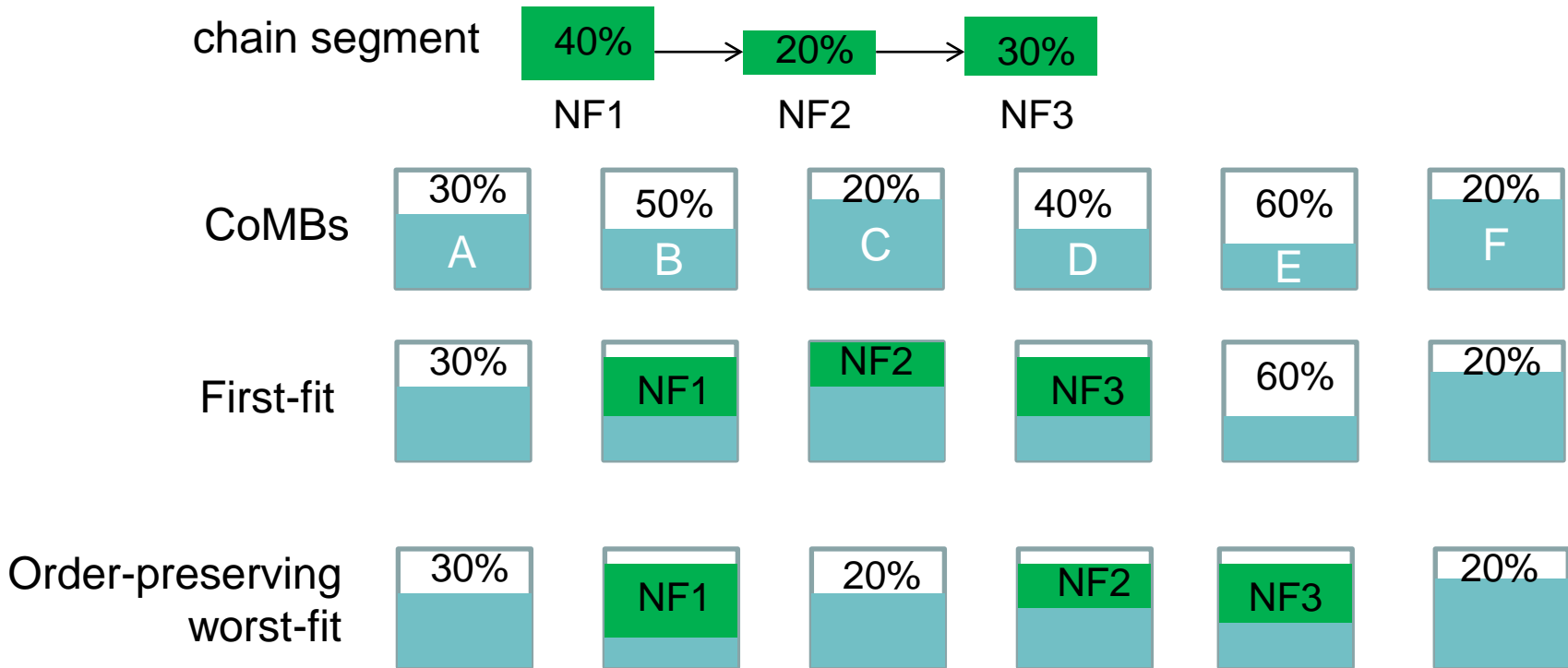


Objectives:

- Load balancing
- Correctness

Approach:

- Step 1: Order-preserving First-fit
- Step 2: Order-preserving worst-fit

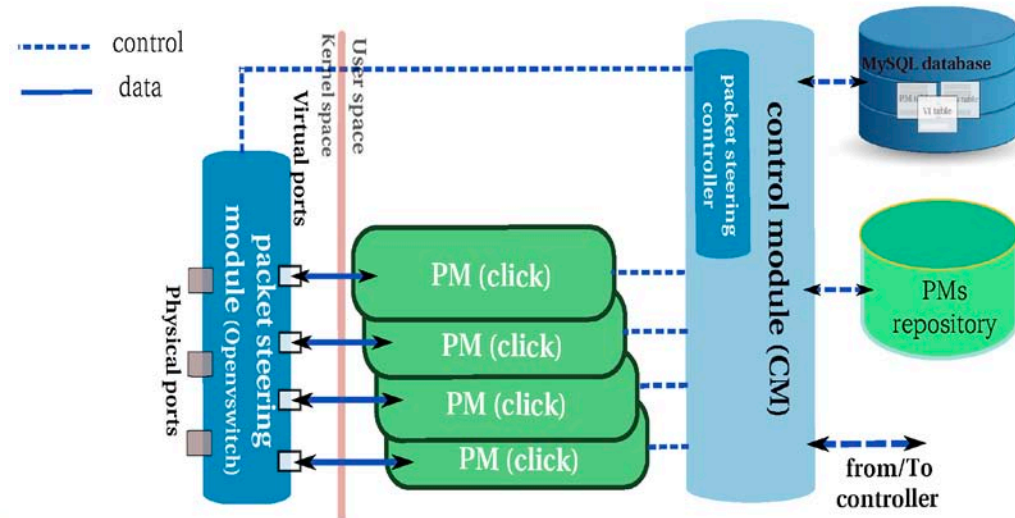




Implementation



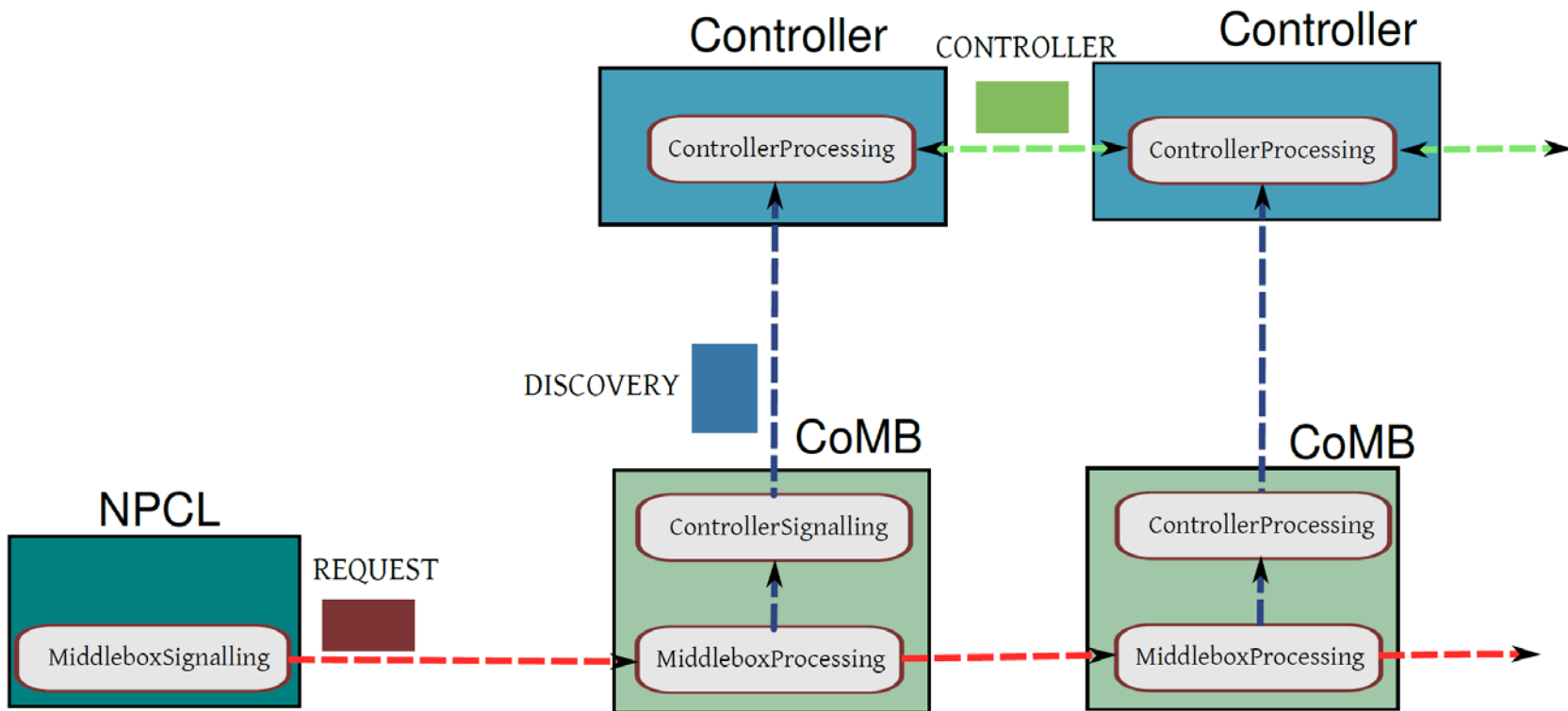
- Processing module (PM):
 - Implements NFs using Click Modular Router
- Packet steering module:
 - Steers traffic between PMs and physical ports using OpenvSwitch
- Repository:
 - Stores PM configuration templates
- Control module:
 - Installs, configures, and terminates PMs
 - API exposed to controller



Discovery implementation



- Implemented 4 Click modular router elements
- Encapsulate messages in UDP packets
- Identify signaling message by the IP Router Alert Option (RAO)

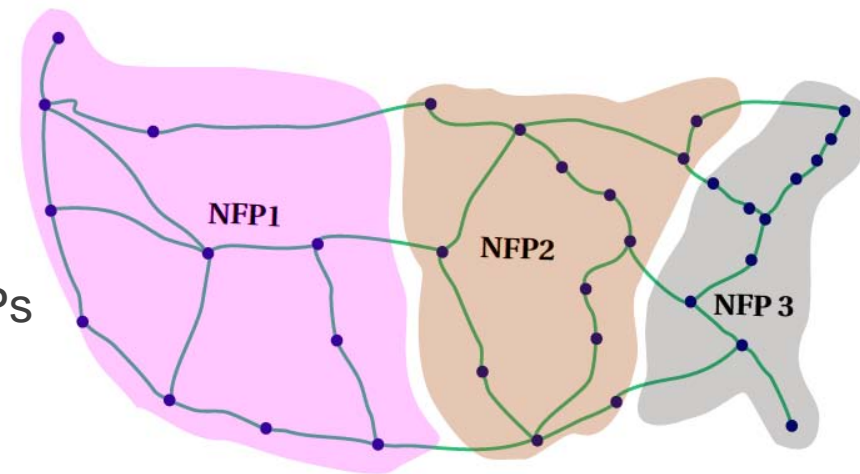




Evaluation

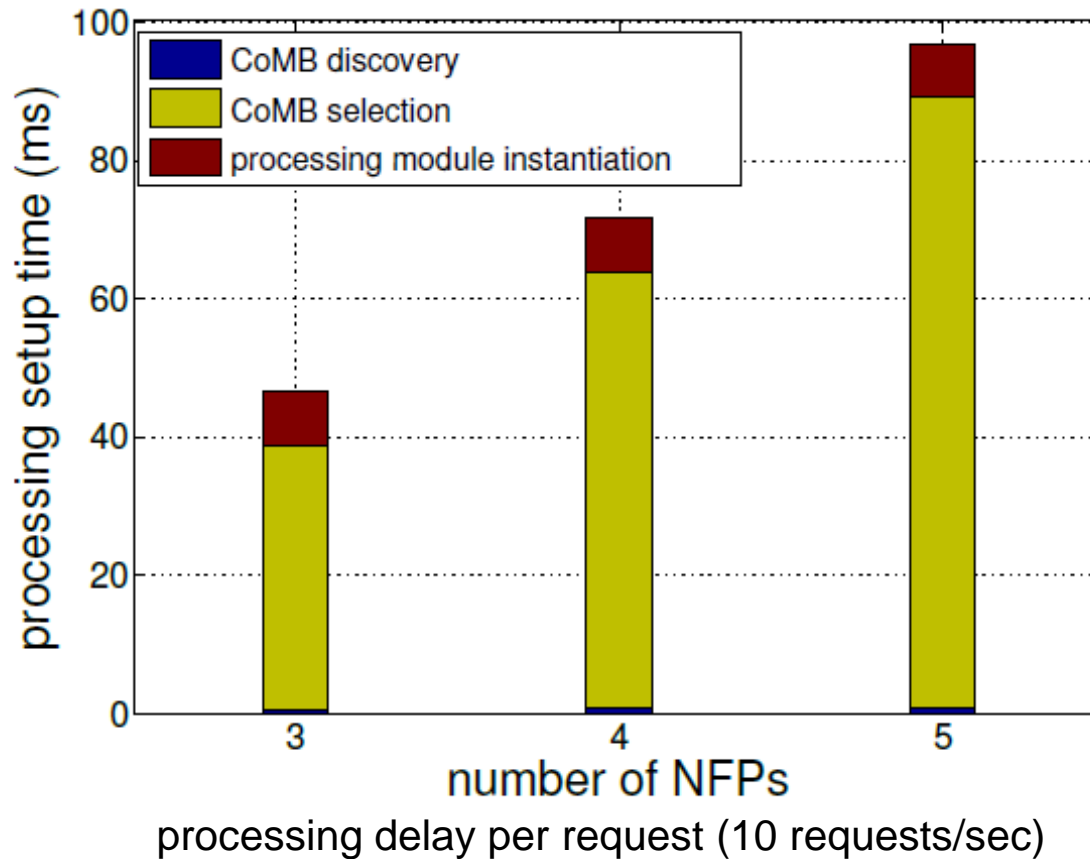


- Experimental evaluation of flow processing setup delay:
 - 22 servers deployed in an Emulab-based testbed (FILAB):
 - quad-core Xeon CPUs @2.27GHz and 6 GB DDR3
 - 2 - 5 NFPs, each with:
 - 1 controller
 - 3 CoMBs (deployed in separate nodes)
- Evaluation of ComB selection efficiency with simulations:
 - Simulator:
 - Flow-level simulator (Python)
 - Simulation setup:
 - Internet-2 topology
 - 34 CoMBs subdivided into 3 NFPs



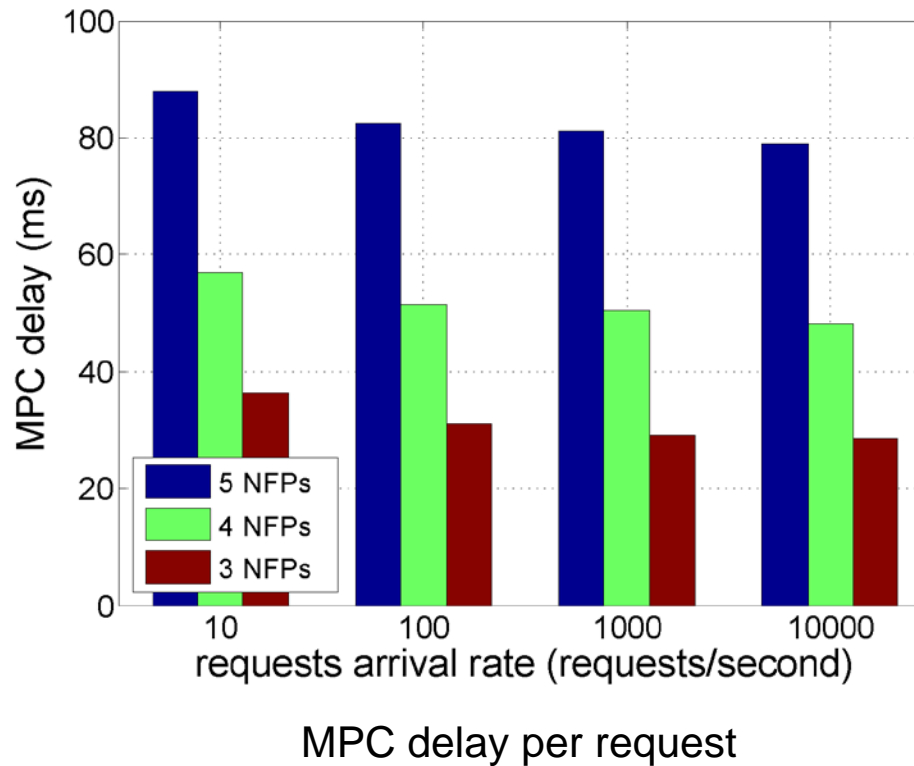


- Flow setup delay < 100 ms
- CoMB selection dominates flow processing setup delay
 - MPC is computationally-intensive ($O(n^2)$, n is the number of NFPs)



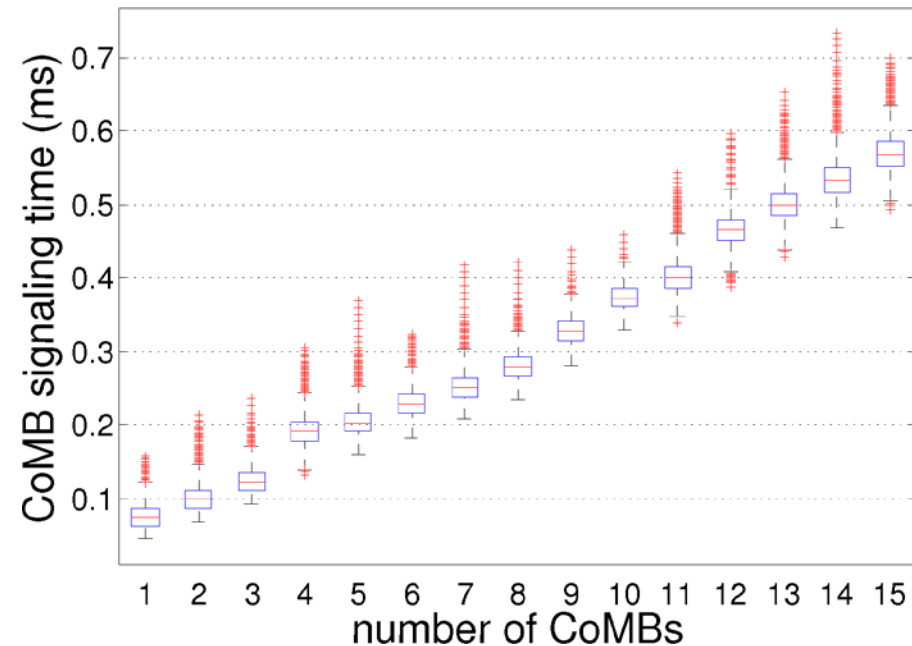


- MPC delay < 100 ms for up to 5 NFPs (i.e., average AS-path length)
 - MPC can be scaled with GPU

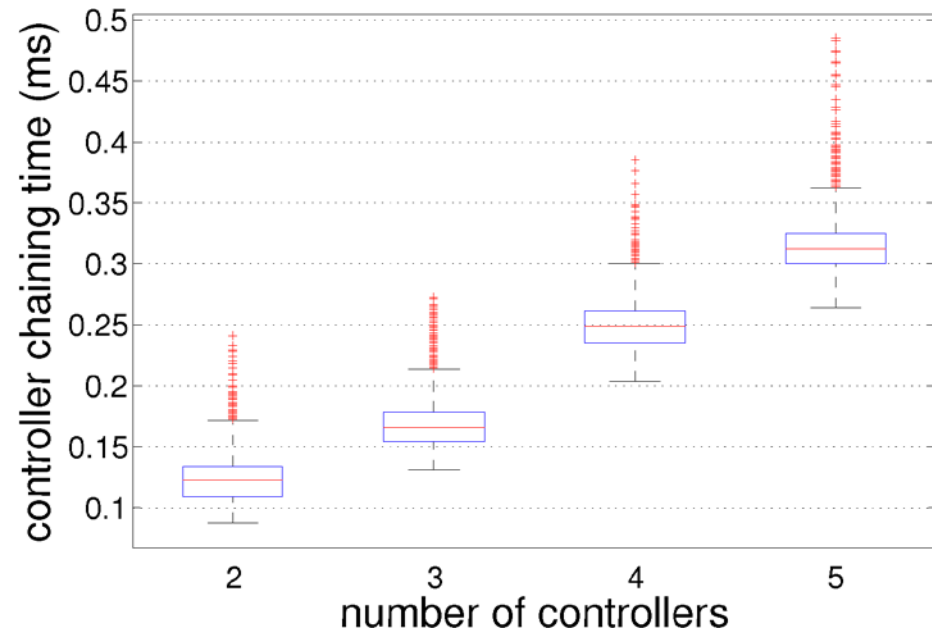




- Minimal delay with CoMB signaling and controller chaining
- Middlebox discovery scales with the number of CoMBs and controllers



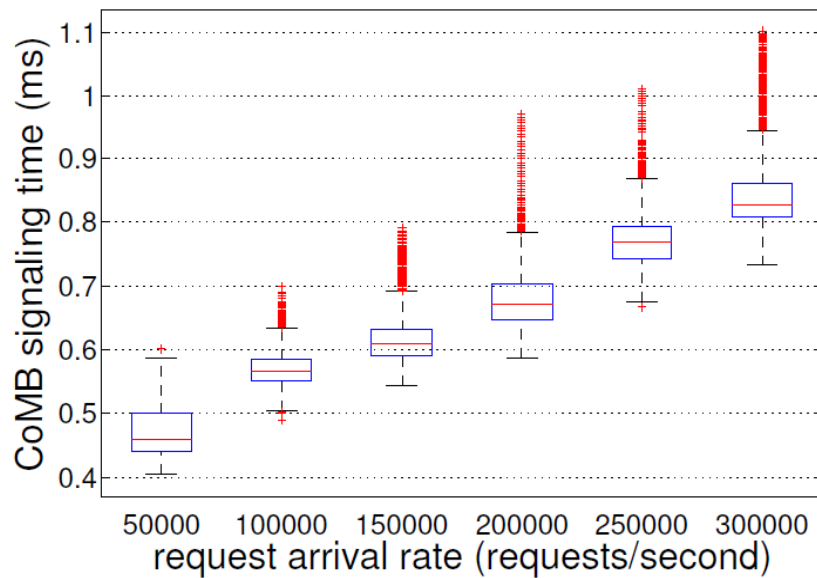
CoMB signaling delay per request



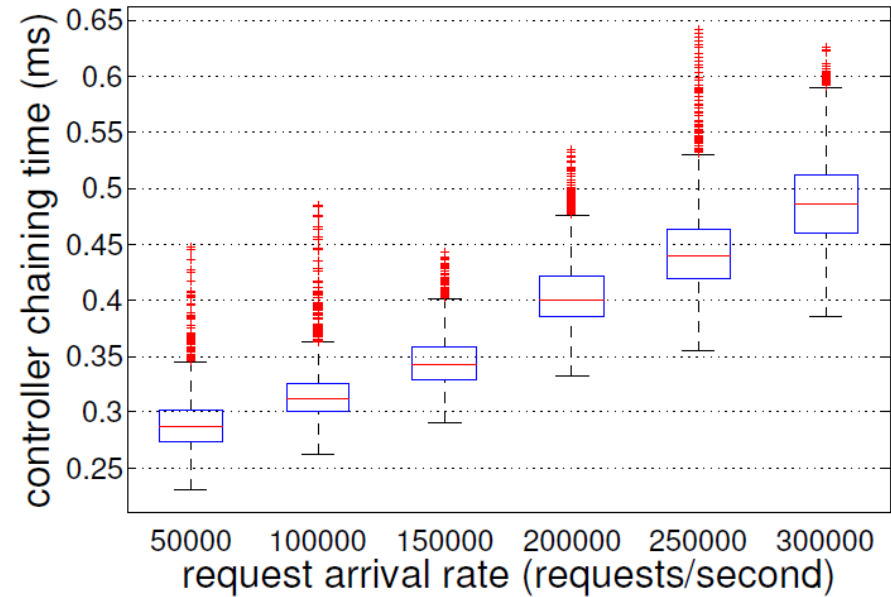
controller chaining delay per request



- Middlebox discovery scales with the number of requests (300K requests/second)



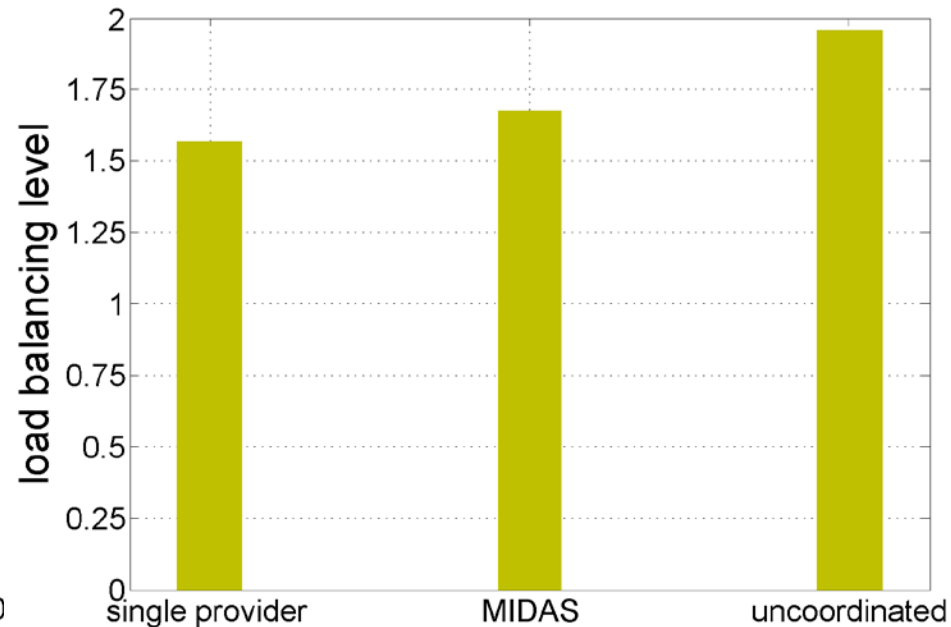
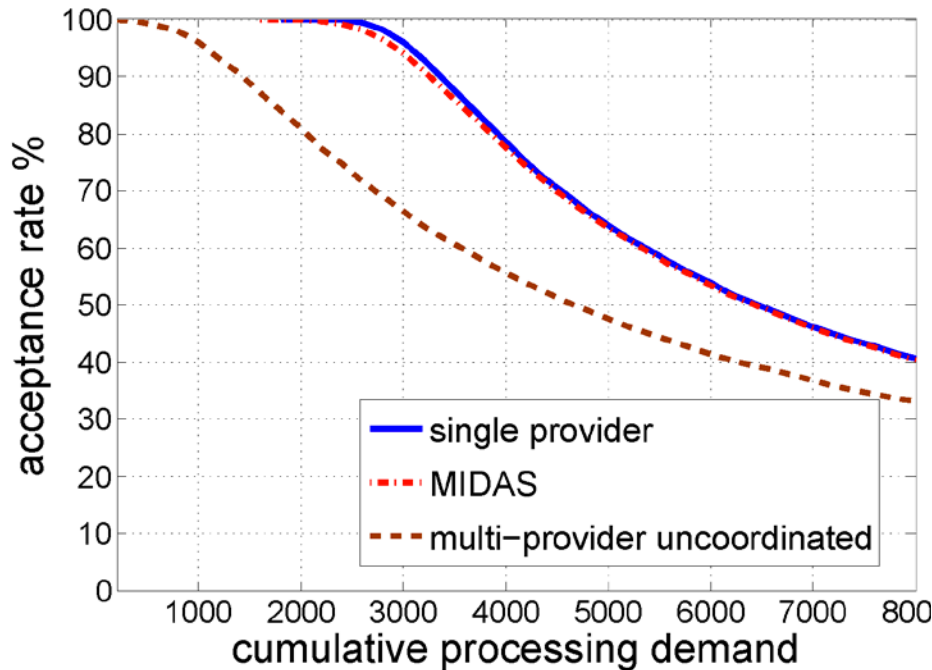
CoMB signaling delay per request



controller chaining delay per request



- Comparison method:
 - Single provider:
 - All CoMBs managed by a single controller
 - Multi-provider uncoordinated:
 - On-the-fly selection of CoMBs based on the utilization level





Conclusions



- MIDAS enables:
 - Middlebox discovery without prior knowledge of the traffic path
 - Privacy-preserving Interoperability among NFPs for middlebox selection
 - Rapid and order-preserving network service embedding
 - Feasibility of coordinated on-path processing setup



Thank you!

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